

CYCLIC SQUARE ROOT OF GRAPHS

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Abstract

A graph G is said to have a cyclic square root if there exists a graph H which is cyclic and is such that $H^2 = G$ (upto isomorphism). A set of necessary and sufficient conditions for the existence of cyclic square roots for graphs is given. These conditions are in terms of cliques. An applications to 4-regular graphs is given.

Introduction

By a graph we mean a connected graph without loops and multiple edges. We follow the notations and terminology of Harary (1). The square of a graph G , written as G^2 , is obtained by adding to G edges which join pairs of vertices of G at a distance 2 apart. A graph H is to be a square root of a graph G if $H^2 = G$ (up to isomorphism). Mukhopadhyay [3] presented a solution to the problem of characterising graphs having at least one square root graph. In this paper we obtain a set of necessary and sufficient conditions for the existence of a square root which is a cycle. We call such square roots as cyclic square roots..;

If G is a graph with at least three vertices, then a clique of G . A 3-clique is a clique on three vertices. If X is a finite set, then $|X|$ will denote the number of elements of X . If a is a vertex of a graph G then the neighbourhood V_a of a is $\{a\} \cup \{b \in V(G) \mid ab \in E(G)\}$. If S is a non-empty subset of $V(G)$ then the induced subgraph of G by S , written $G \langle S \rangle$, is the maximal subgraph of G with vertex set S . If we square a 3-cycle we get a 3-cycle and squaring the 4-cycles and 5-cycles yield the complete graphs on 4 and 5 vertices respectively. We therefore consider graphs with atleast 6 vertices.

Proposition : Let G be a non-complete graph which is the square of some cycle H , where $|V(H)| = |V(G)| \geq 6$. for

$a \in V(H) = V(G)$, let V_a denote the neighbourhood of a in H and $K_a = G \langle V_a \rangle$; Then the following hold.

(i) K_a is a 3- clique in G , for each $a \in V(H)$

(ii) $V(K_a) \cap V(K_b) = \{a, b\}$ if and only if a is adjacent to b in H .

(iii) a is adjacent to c and c is adjacent to b in H if and only if
 $V(K_a) \cap V(K_b) = \{c\}$

(iv) $V(K_a) \cap V(K_b) = \emptyset$ if and only if every path from a to b in H is of length ≥ 3 .

(v) For each 3- clique in G there exist exactly two other 3- cliques each intersecting with it in exactly two vertices of G .

Proof :

(i) Let $u, v \in V(K_a) = V_a$. The distance $d(u, v) \leq 2$ in H . Therefore $uv \in E(H) = E(G)$ and $u, v \in V(K_a)$. Hence K_a is complete. Now $a \in V(H)$ and H is a cycle. Therefore $|V(K_a)| = 3$. That K_a is maximal, follows from the fact that $|v(H)| \geq 6$.

Proofs of (ii), (iii) and (iv) are routine. We prove (v).

For $u_i \in V(H)$, by (i) $K_i = G \langle V_i \rangle$ is a 3- clique where V_i is the neighbourhood of U_i in H . But H is a cycle. Therefore there exist u_j and u_k in $V(H)$ such that u_i is adjacent to u_j and u_k in H . Thus by (ii) $V(K_i) \cap V(K_j)$ contains exactly two vertices and $V(K_i) \cap V(K_k)$ also contains exactly two vertices. Therefore there exist two cliques K_j and K_k which intersect with K_i in exactly two vertices. Now suppose there exist a third clique different from K_j and K_k , say K_r such that $V(K_i) \cap V(K_r)$ contains exactly two vertices. Then u_i is also adjacent to u_r different from u_j and u_k , that is, the degree of u_i is 3 or more in H , which is a contradiction, since H is a cycle. Therefore K_i meets two cliques in exactly two vertices.

Theorem 1: A non-complete graph G on p vertices u_1, u_2, \dots, u_p has a cyclic

square root H if and only if there exists a collection of p 3-cliques K_1, K_2, \dots, K_p such that:

- (a) $U_i \in V(K_i)$ for every i .
- (b) $\bigcup_{i=1}^p E(K_i) = E(G)$.
- (c) No two K_i s intersect in more than two vertices.
- (d) For each K_i there exist exactly two other K_j and K_k such that K_i meets each of K_j and K_k in exactly two vertices.
- (e) $u_i \in V(K_j)$ if and only if $u_j \in V(K_i)$ for every i and j .

Proof : Suppose G has a cyclic square root H . For each $u_i \in V(H)$, let V_i be the neighbourhood of u_i in H and $K_i = G \langle V_i \rangle$. From proposition (i) each K_i is a 3-clique in G . Thus we have a collection of p 3-cliques in G . Now we have to show that at conditions (a) -(e) are satisfied. (a) is immediate from the definition of V_i (e) follows the fact that u_i is in the neighbourhood of u_j if and only if u_j is in the neighbourhood of u_i in H . (c) and (d) follow from (ii), (iii), (iv) and (v) of proposition.

To prove (b) : let $u_i, u_j \in E(G) = E(H^2)$. Therefore $d(u_i, u_j) \leq 2$ in H . Hence there exists $u_k \in V(H)$ such that u_i, u_k, u_j is a path in H . That is $u_i, u_j \in E(K_k)$. Hence

$$E(G) \subseteq \bigcup_{i=1}^p E(K_i). \text{ The reverse inclusion is obvious.}$$

Conversely, we define the graph H as.

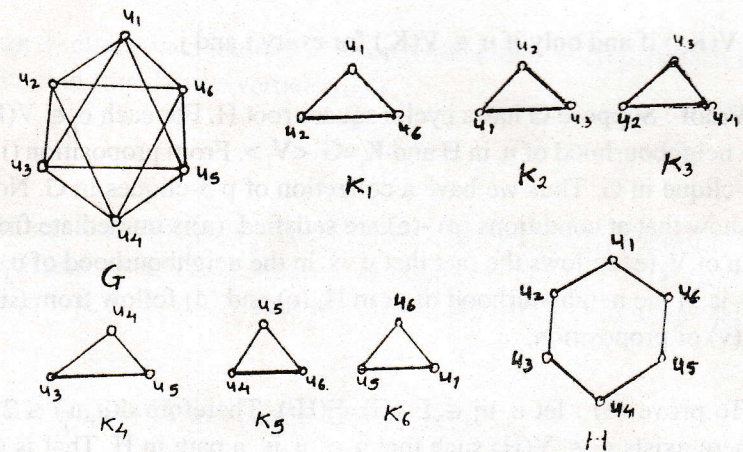
$V(H) = V(G) = \{u_1, u_2, \dots, u_p\}$ and
 $E(H) = \{u_i, u_j \in E(G) / V(K_i) \cap V(K_j) = \{u_i, u_j\}\}$ We show that
 $E(H^2) = E(G)$. Let $u_i, u_j \in E(H^2)$. If $d(u_i, u_j) = 1$ in H then $u_i, u_j \in E(H) \subseteq E(G)$.
 Otherwise, there exists $u_k \in V(H)$ such that u_i, u_k, u_j is a path in H . Therefore $u_i, u_j \in V(K_k)$ and since K_k is complete, $u_i, u_j \in E(K_k)$ for some k . By (b) $u_i, u_j \in$

$E(G)$. Thus $E(H^2) \leq E(G)$. Again let $u_i, u_j \in E(G)$. From (b), (c) and (e) $V(K_i) \cap V(K_j) = \{u_i, u_k\}$ and $V(K_j) \cap V(K_k) = \{u_j, u_k\}$. Hence $u_i, u_k \in E(H)$ and $u_j, u_k \in E(H)$. That is u_i, u_k, u_j is a path in H . Therefore $u_i, u_j \in E(H^2)$, which proves that

$E(G) \leq E(H^2)$. Thus $E(G) = E(H^2)$. Also $|V(H)| = p$ and clearly H is connected. By (d) degree of each vertex of H is two. Hence H is a cycle and $H^2 = G$.

Remark. A procedure for finding all the cliques in a graph is known [2]. As a result, our Theorem gives a procedure for finding a cyclic square root of a 4-regular graph in which every clique is a 3-clique.

We illustrate the Theorem by an example.



Figure

A graph G and all its 3-cliques $k_1 - k_6$ are shown in figure. It is easily seen that the conditions of Theorem are satisfied by k_1 to k_6 . The cycle H is a cyclic square root of G .

Application:

Following theorem is an application of our Theorem 1.

Theorem 2: Any two 4-regular graphs with the same number of vertices having cyclic square roots are isomorphic.

Proof : Let G_1 be a 4- regular graph on p vertices u_1, u_2, \dots, u_p having cyclic square root c_1 , say G_2 be a 4- regular graph on p vertices v_1, v_2, \dots, v_p having cyclic square root c_2 , say. We can assume without loss that mapping $f : V(C_1) \rightarrow V(C_2)$ defined by $f(u_i) = v_i$ for each i is an isomorphism of c_1 onto c_2 . We show that f preserves the adjacency and non-adjacency as a mapping of G_1 onto G_2 . Let therefore u_i and u_j in G_1 . If u_i and u_j are adjacent in c_1 then $f(u_i)$ and $f(u_j)$ are obviously adjacent in G_2 . Suppose $d(u_i, u_j) = 2$ in C_1 . A path of length 2 in c_1 is carried to a path of length 2 in C_2 by f . Hence $f(u_i)$ and $f(u_j)$ are adjacent in G_2 . By symmetry, f preserves non-adjacency as well.

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